

A Genetic Algorithm for Goal-Conflict Identification

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*in collaboration with
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Early phase in the RE process

Domain Properties

LTL
formulation

Goals

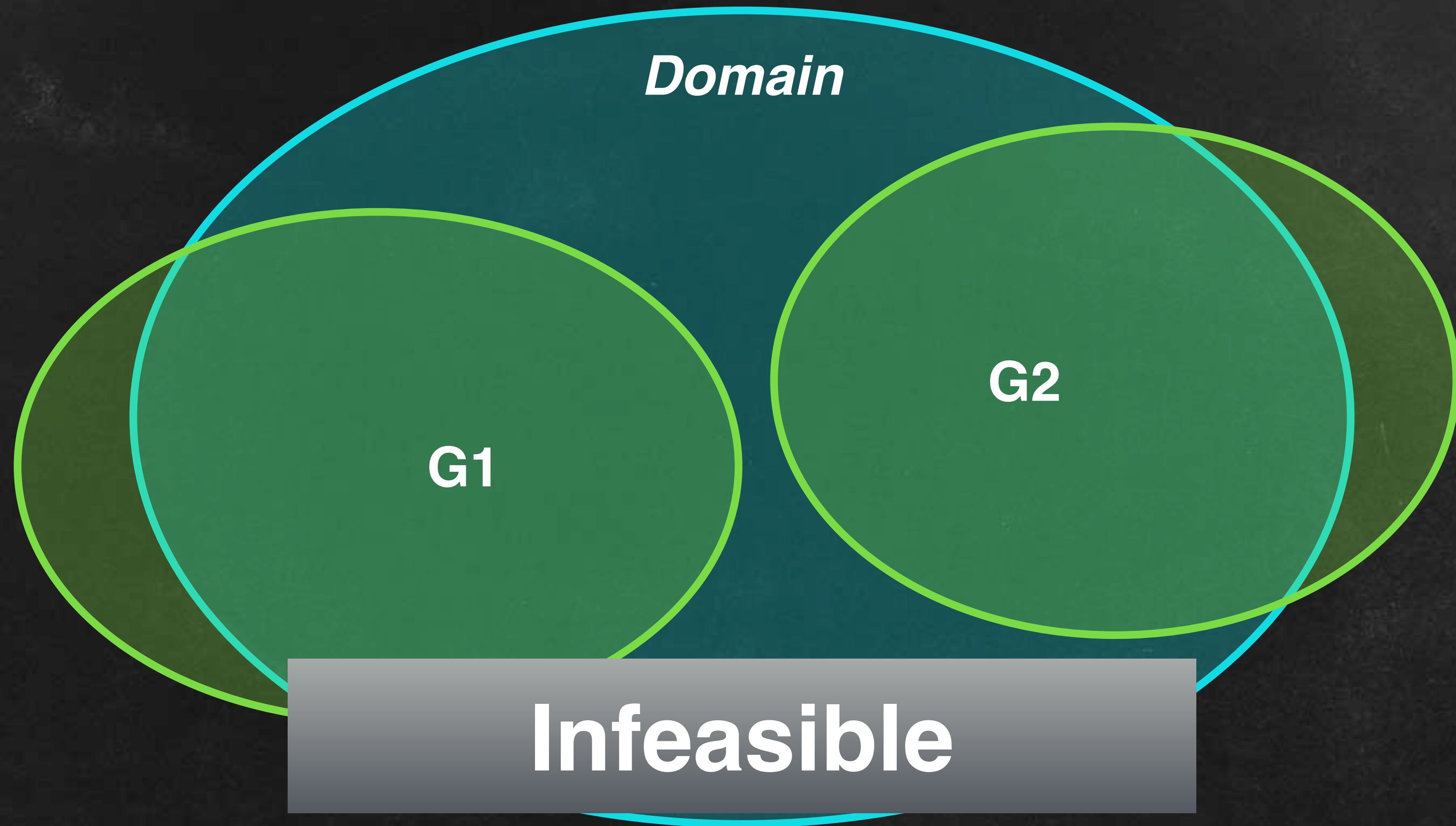
Goals

Goals

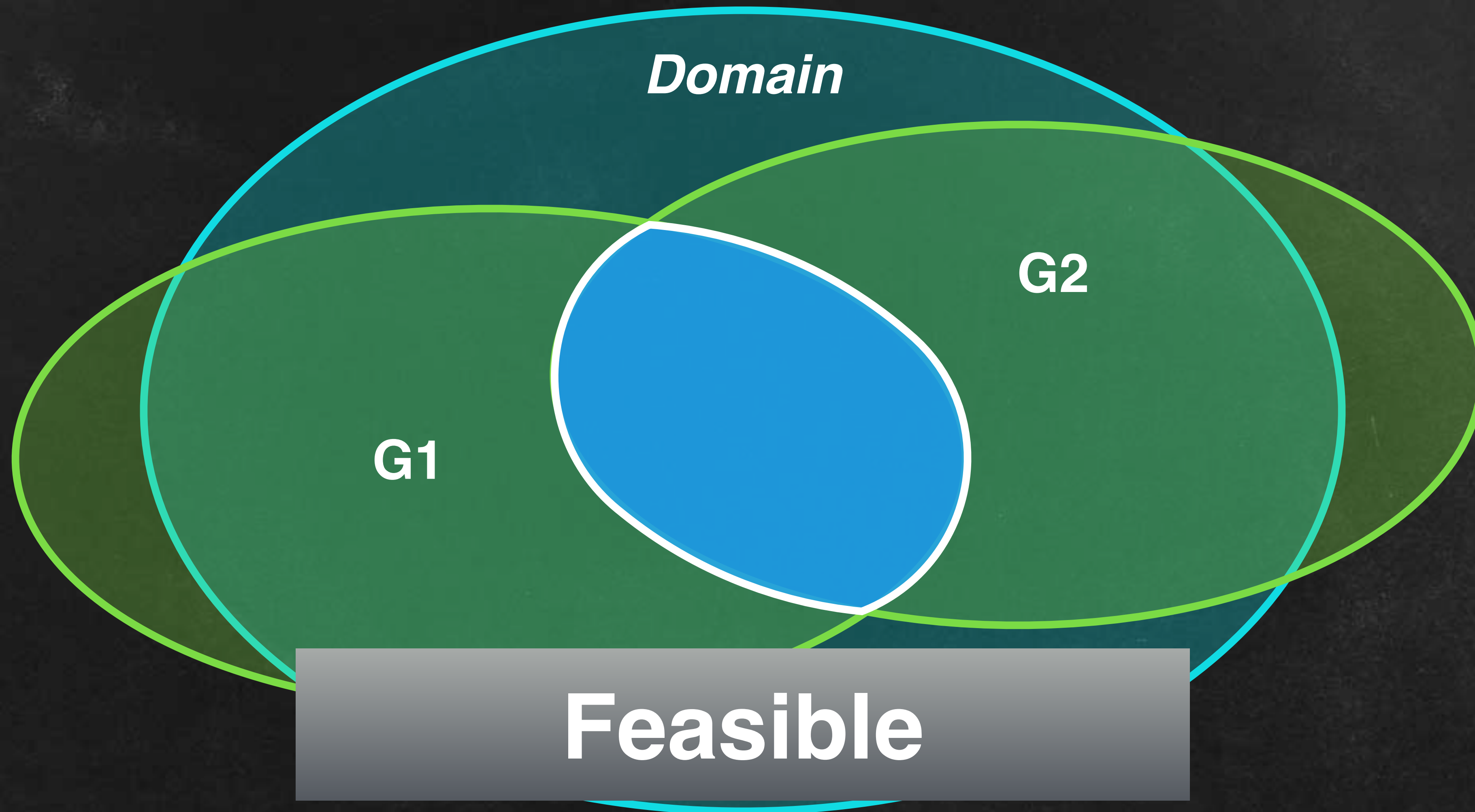
Stakeholders



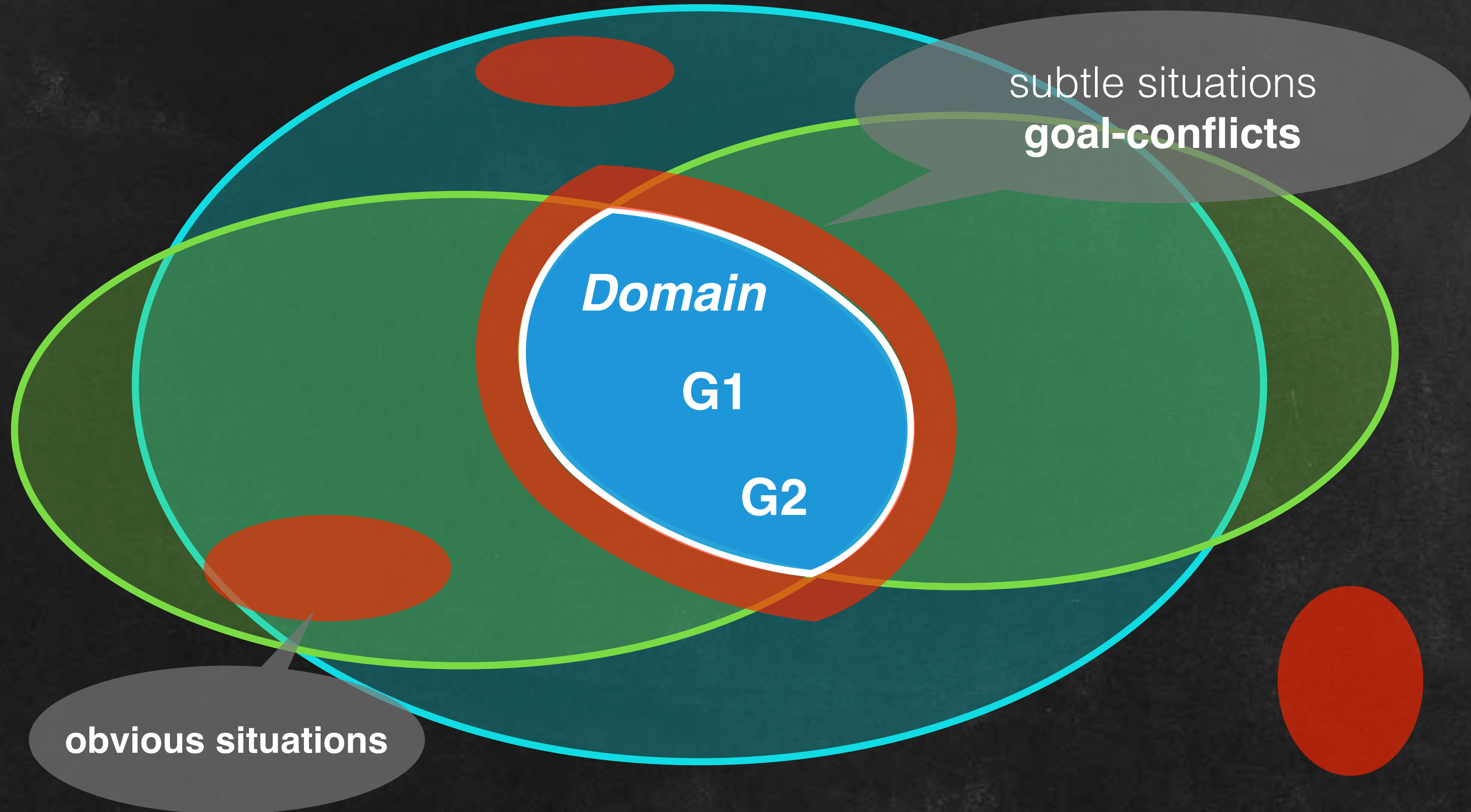
Feasibility



Feasibility

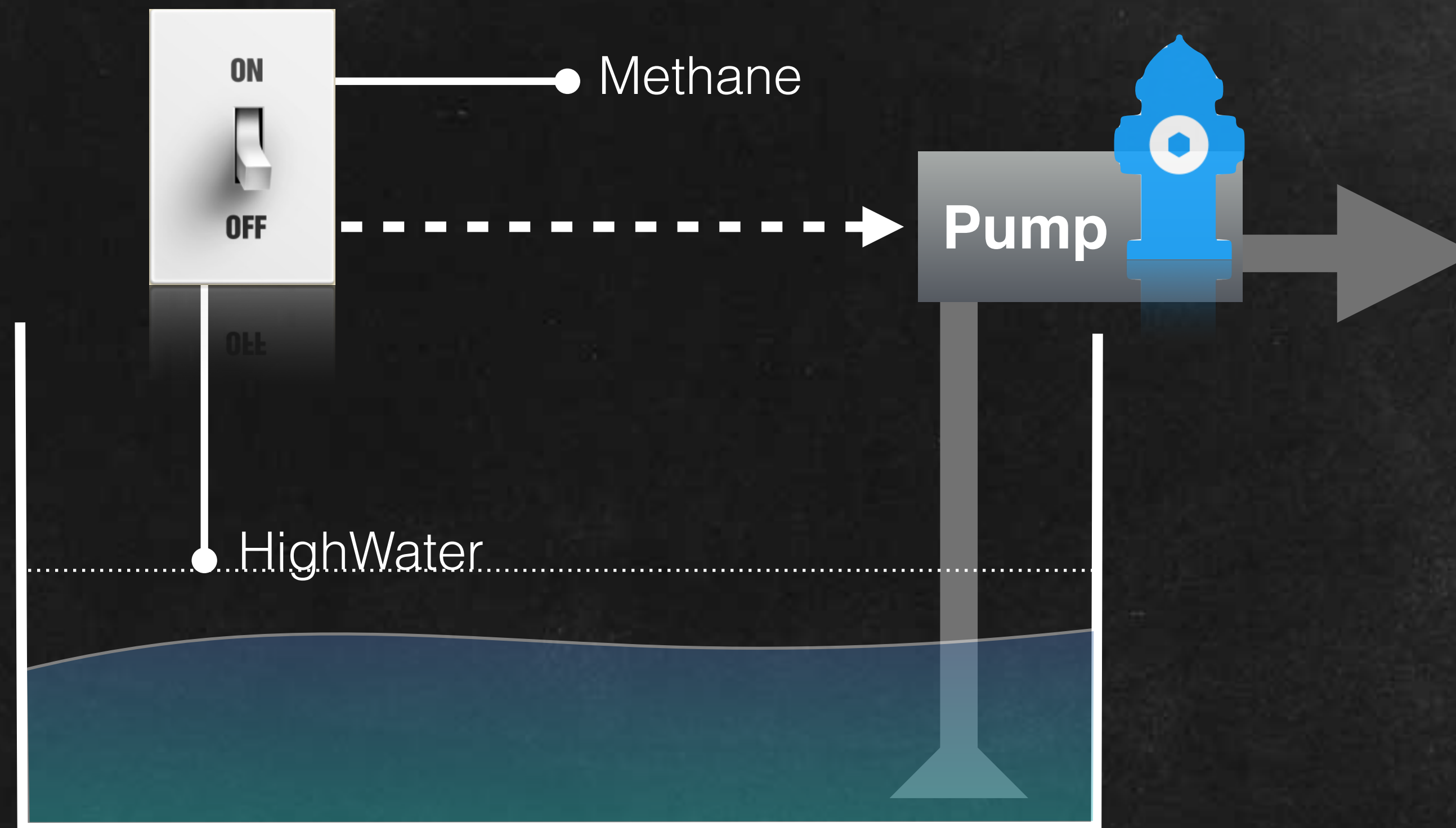


In which situations the goals
can't be fulfilled?



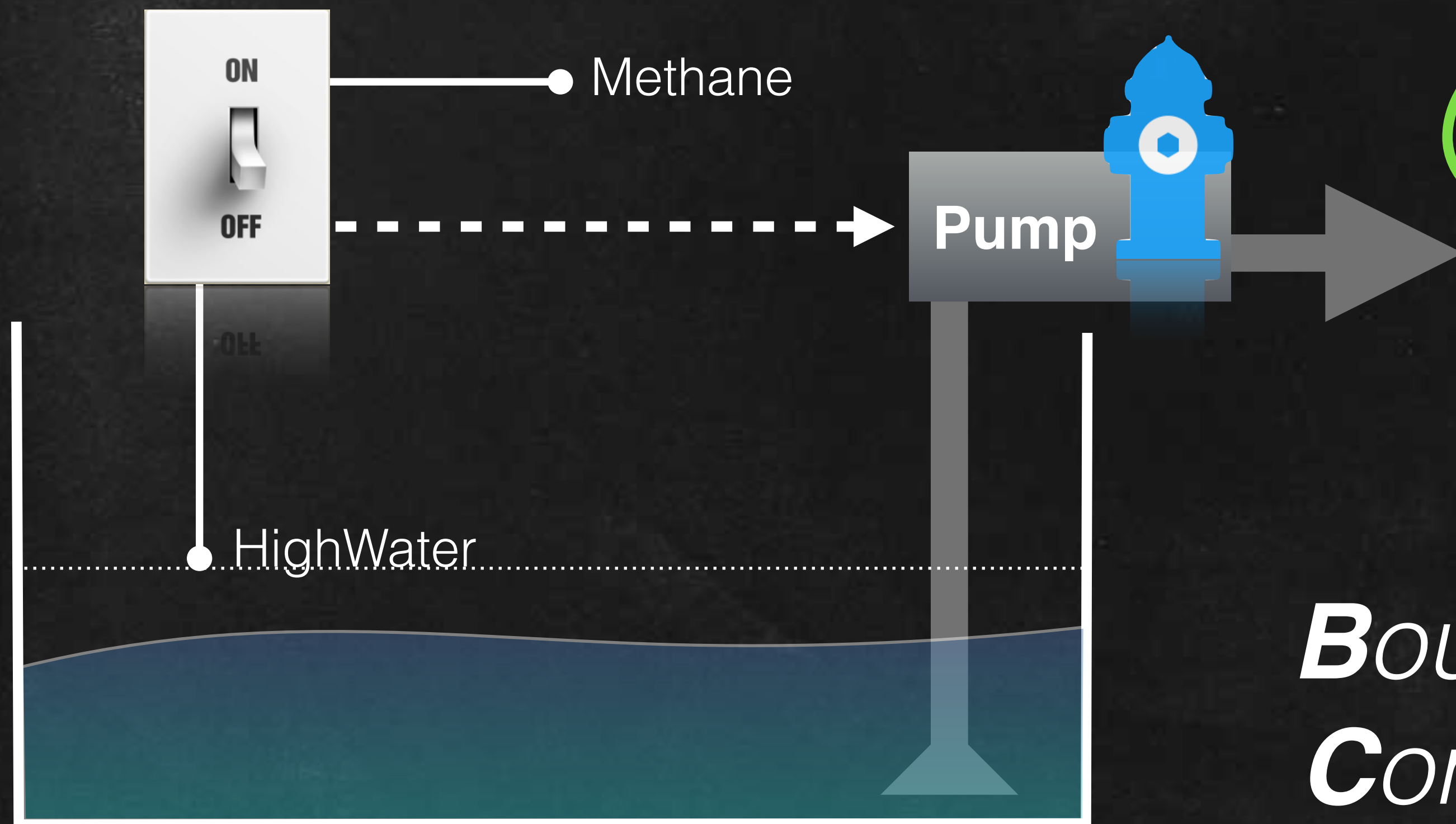
Mine Pump Controller

Mine Pump
Controller



Mine Pump Controller

Mine Pump Controller



If **PumpOn**, then **not HighWater** in at most two minutes

If **Methane**, then **not PumpOn**

If **HighWater**, then **PumpOn**

Boundary Condition

Methane and HighWater



Divergences



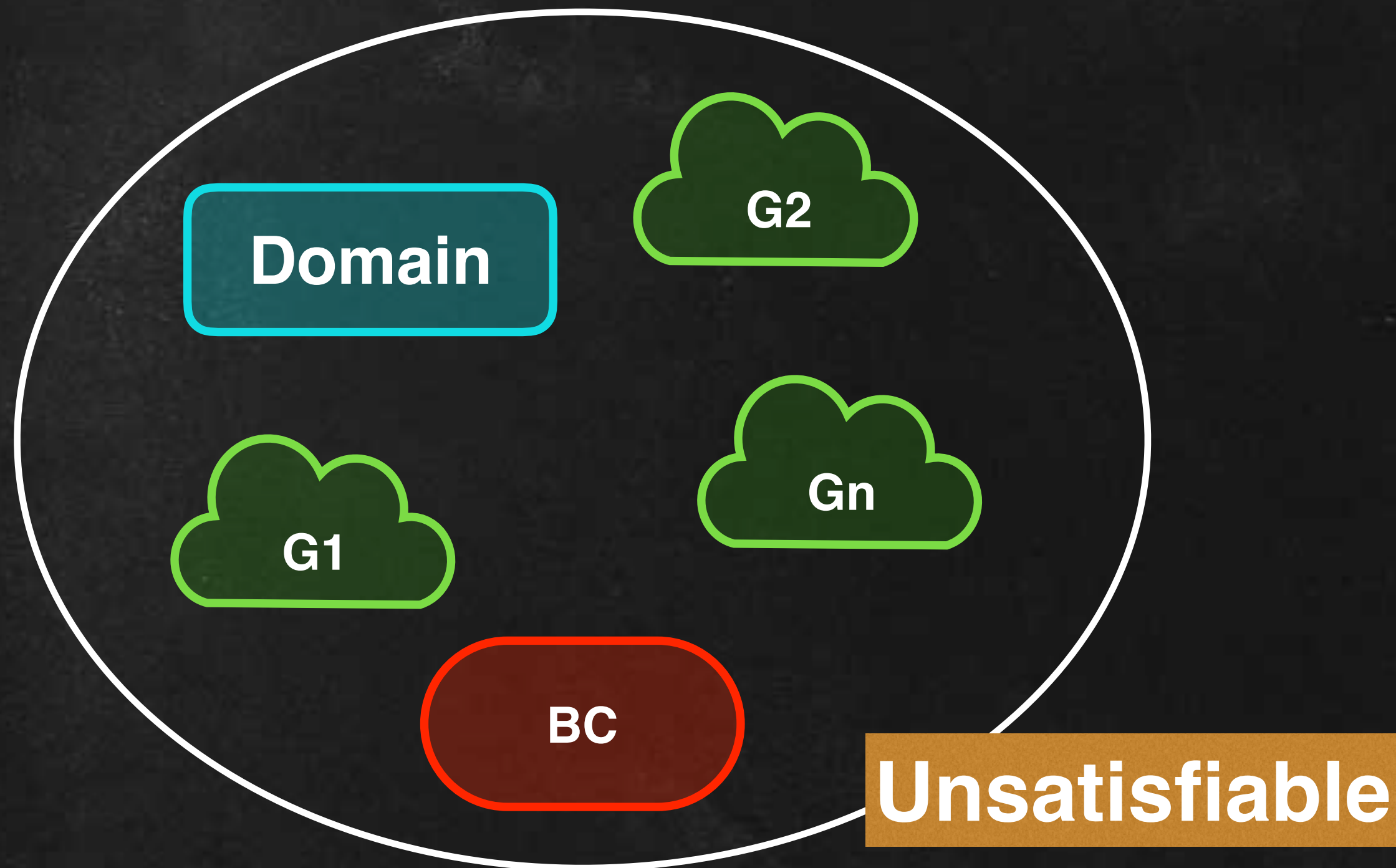
the goals are **divergent** w.r.t. the domain iff there exists a **boundary condition *BC*** such that:

- (1) logical inconsistency
- (2) minimality
- (3) non-triviality

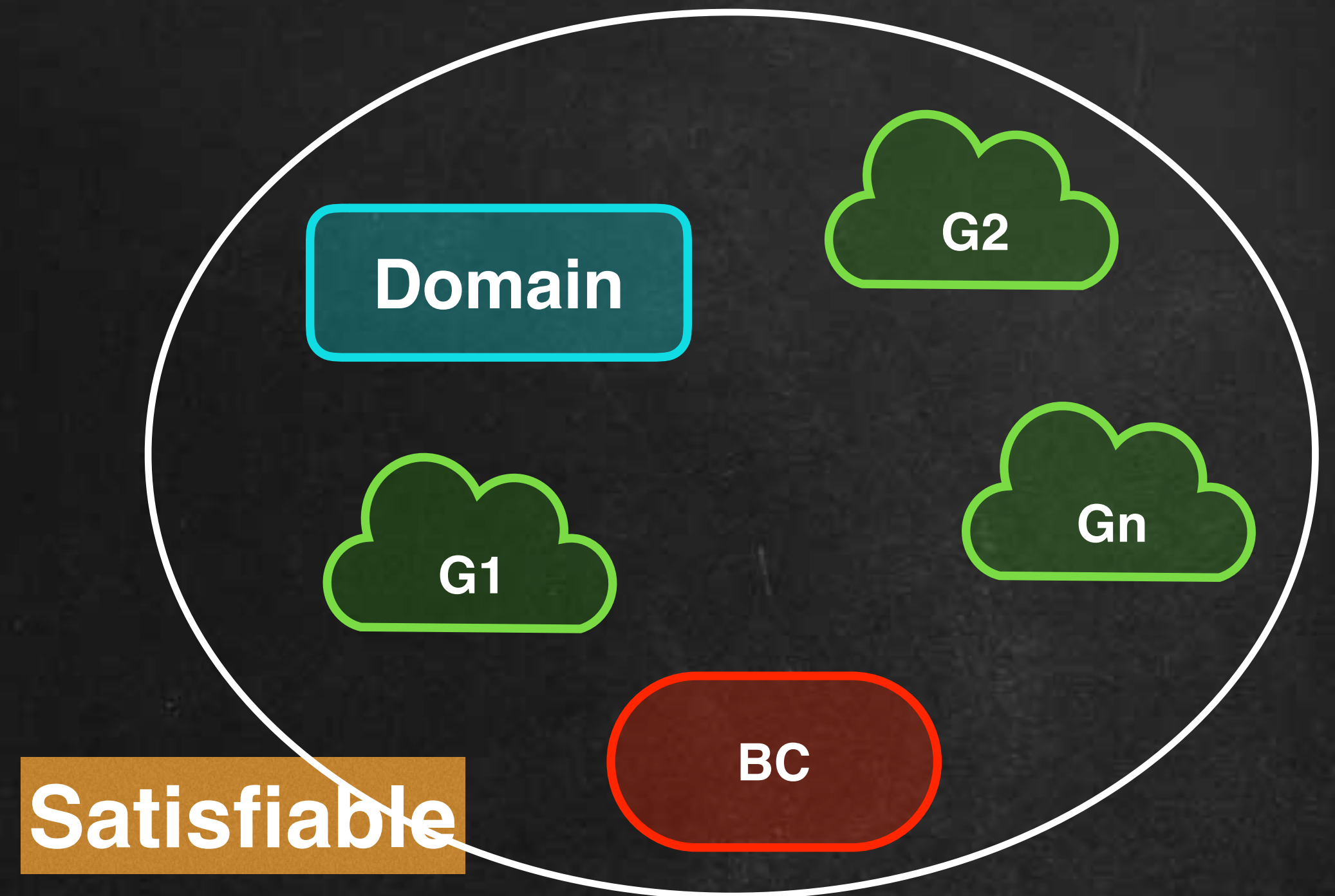


Divergencies

(1) logical inconsistency



(2) minimality

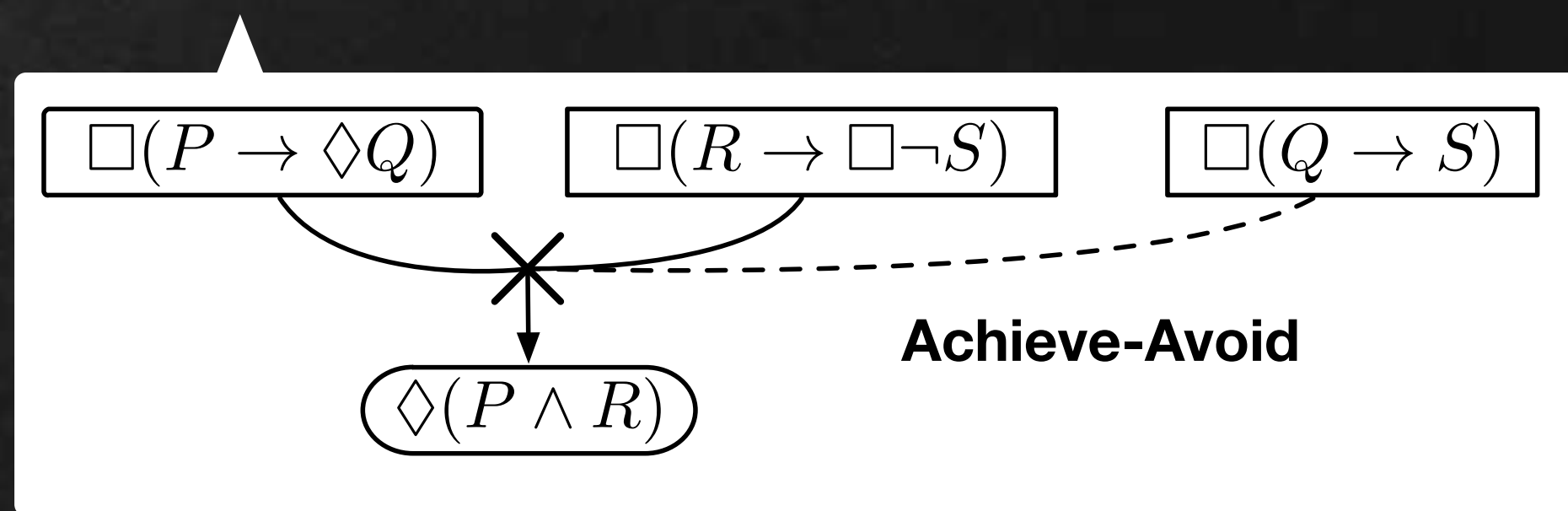


(3) non-triviality $BC \neq \neg(G_1 \wedge \dots \wedge G_n)$

State of the Art

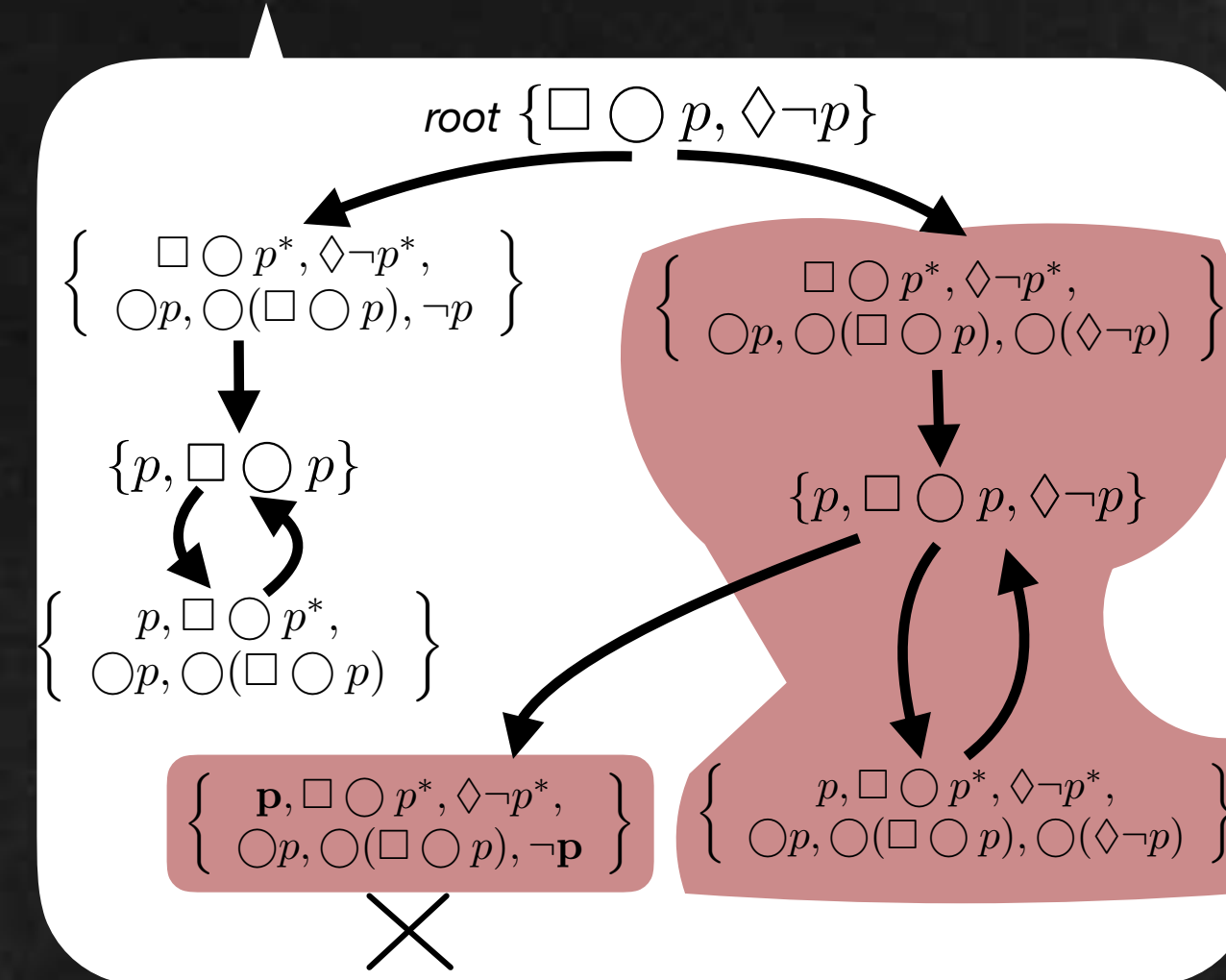
- Automatically identifying boundary conditions

Pattern based technique [TSE'98].
 — restricted to captured patterns.

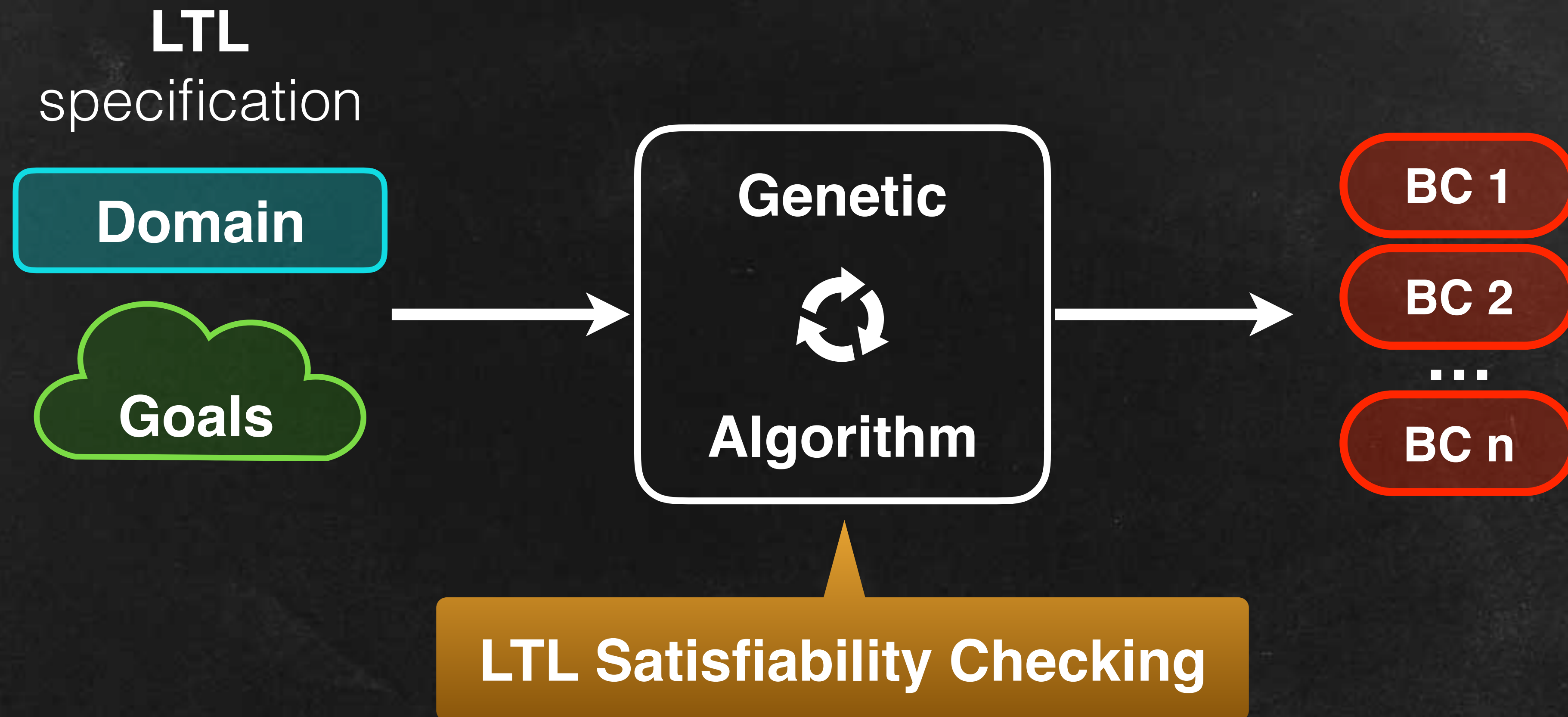


Tableaux based technique [ASE'16]

— very expensive logical manipulation of the tableau structure.



Our Proposal



Genetic Algorithm

Search Space
 $\Box(p \rightarrow \neg q)$ $\Diamond(r \wedge s)$
 $(p \vee q) \mathcal{U} \neg r$

Initial Population

Evaluation

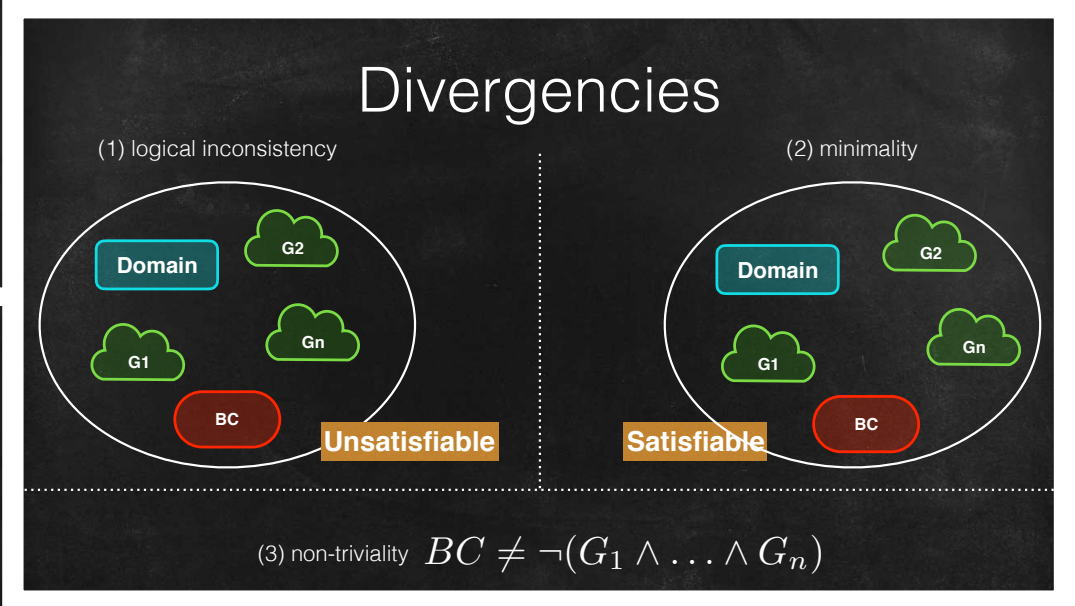
Stop Criterion?

Selection

Genetic Operators

- Mutation
- Crossover

Syntactical alterations
 $\Box(p \rightarrow \neg q)$
 \downarrow
 $\Box(p \wedge \neg q)$



LTL SAT checks

Solutions

Initial Population

- All **sub-formulas**, and their **negations**, computed from the domain properties and the goals.

Goal
 $\Box(p \rightarrow \neg q)$

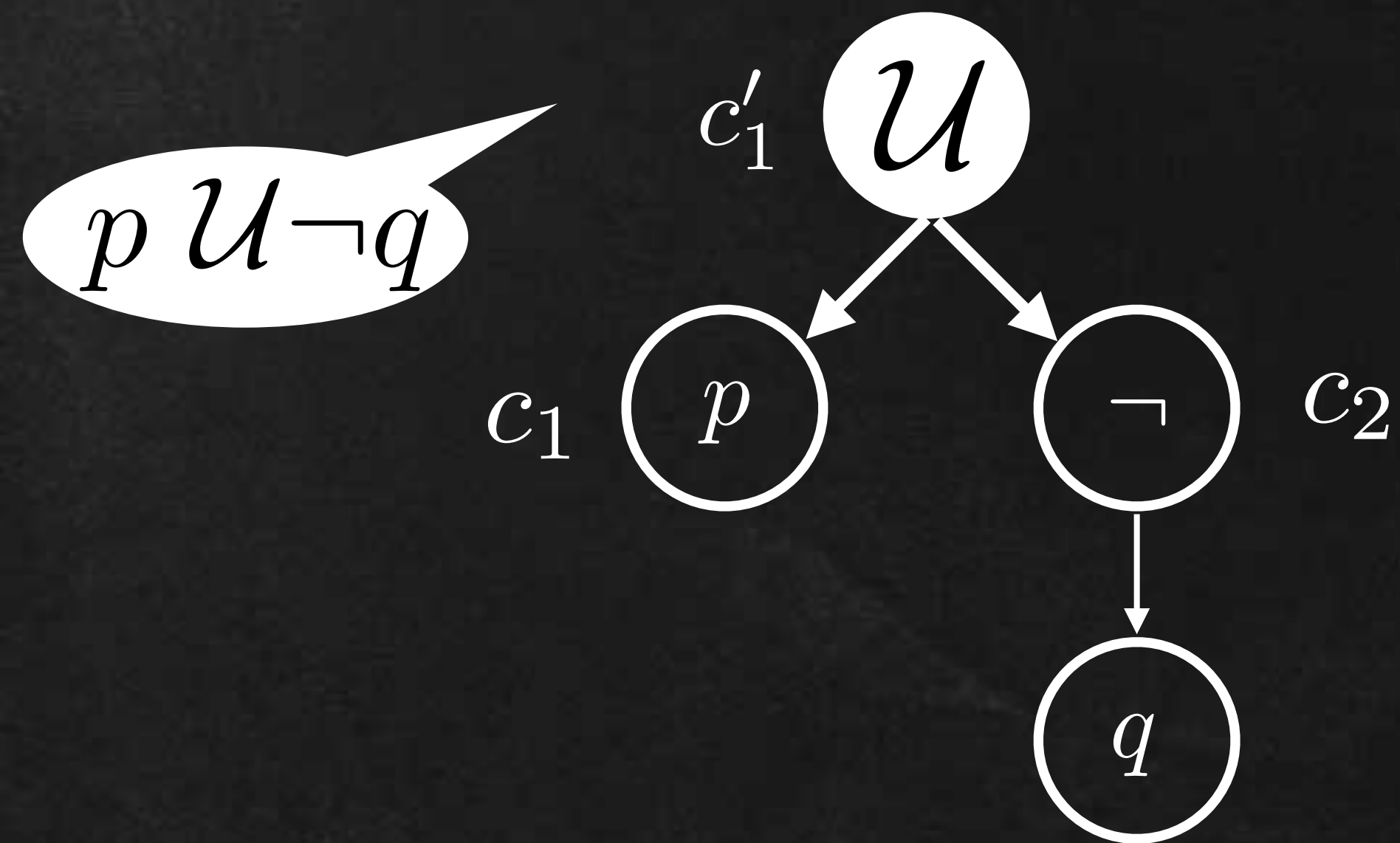
<i>Sub-formulas</i>	<i>Negations</i>
$\Box(p \rightarrow \neg q)$	$\neg\Box(p \rightarrow \neg q)$
$p \rightarrow \neg q$	$\neg(p \rightarrow \neg q)$
$p \quad \neg q \quad q$	$\neg p$

initial population

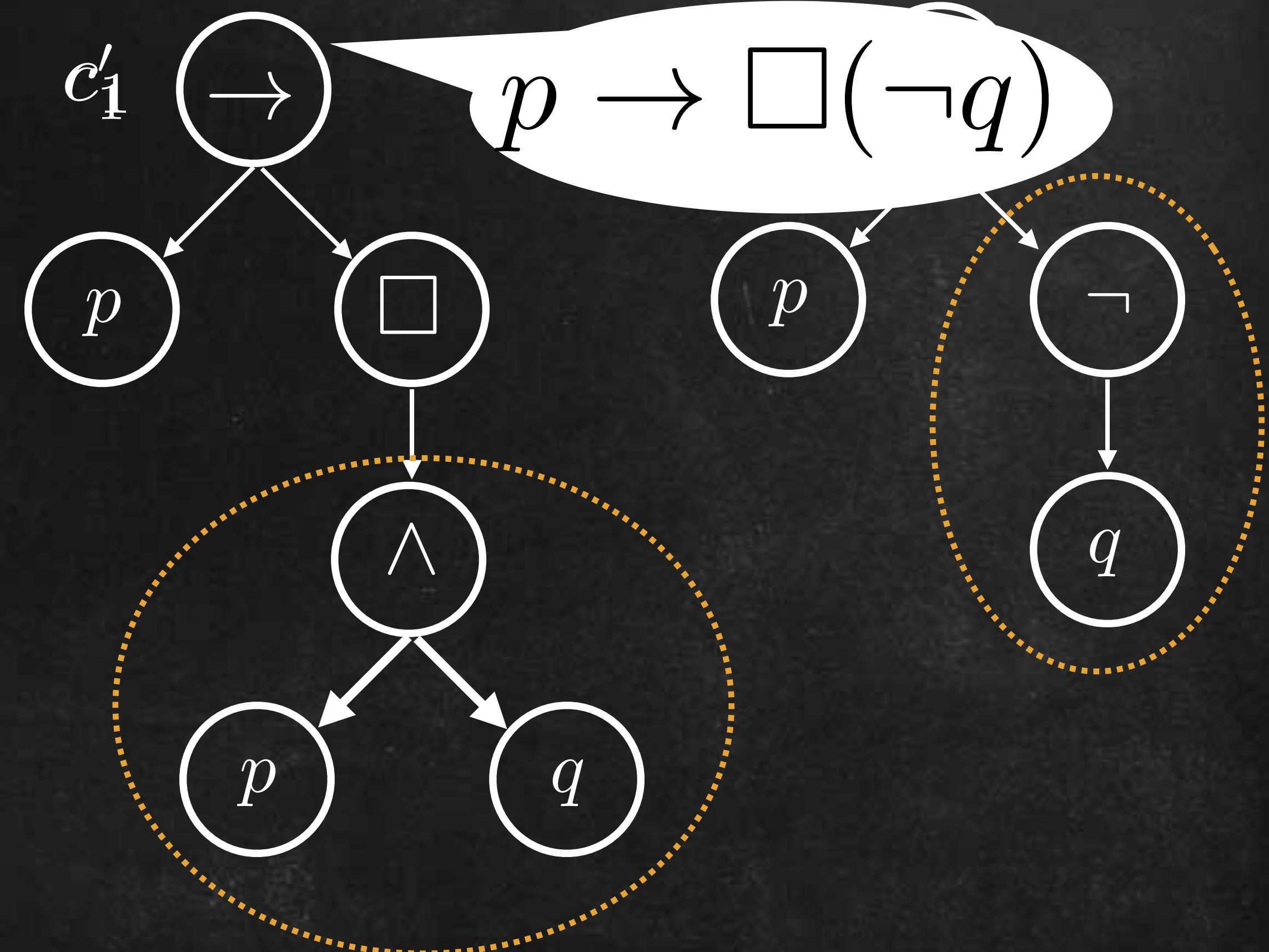
Genetic Operators

Crossover

binary combination



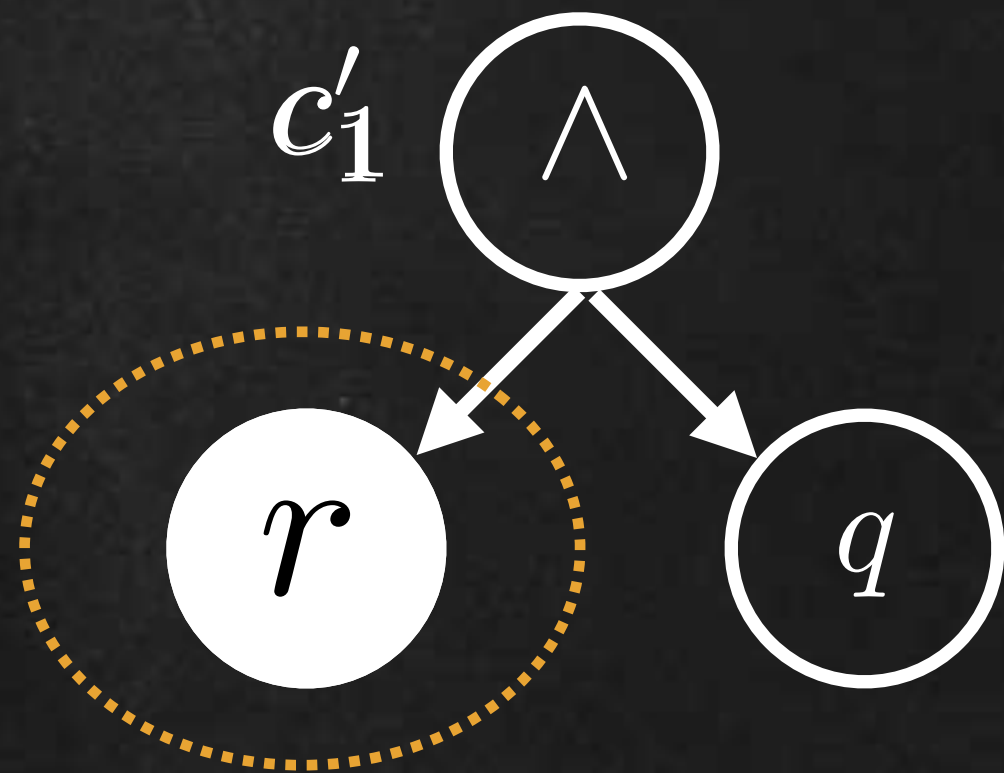
sub formula replacement



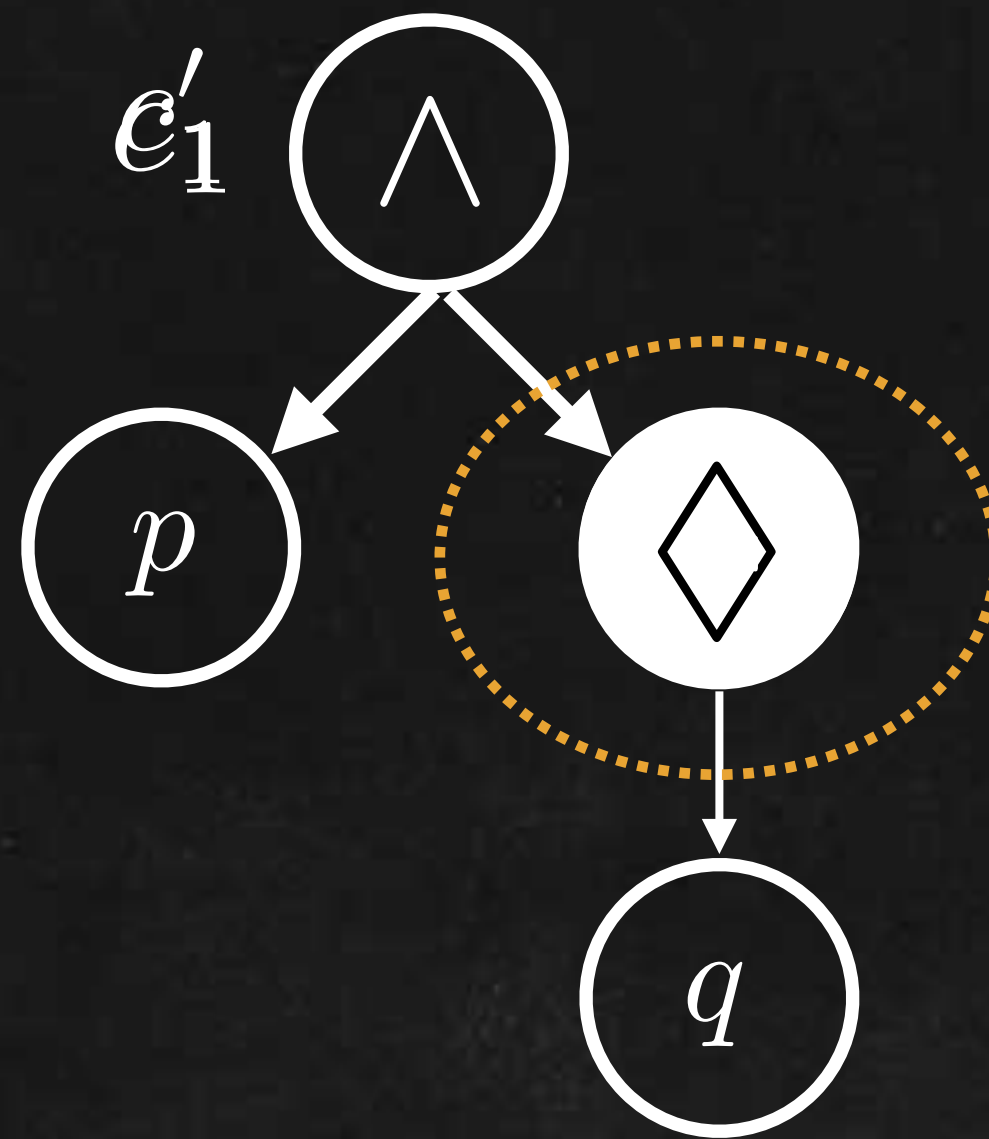
Genetic Operators

Mutation

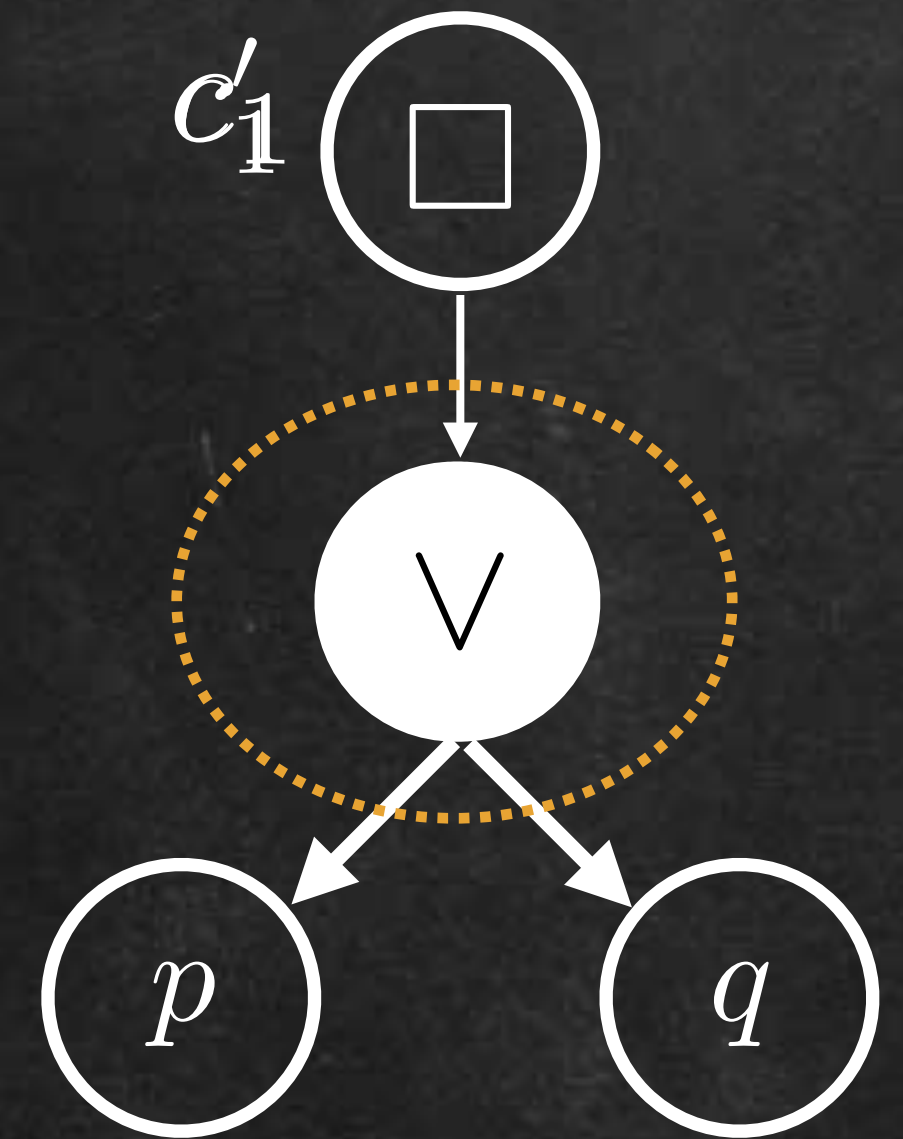
atomic replacement



unary op. replacement



binary op. replacement



Fitness Function

Let φ_c be a candidate **boundary condition**

$$f(\varphi_c) = li(\varphi_c) + \sum_{i=1}^{|G|} \min(\varphi_c, G_i) + nt(\varphi_c) + \frac{1}{\#\varphi_c}$$

logical inconsistency

minimality

non-triviality

formula size
penalty

$$li(\varphi_c) = \begin{cases} 1 & \text{if} \\ 0 & \text{otherwise} \end{cases} \quad \min(\varphi_c, G_i) = \begin{cases} \frac{1}{|G|} & \text{if} \\ 0 & \text{otherwise} \end{cases} \quad nt(\varphi_c) = \begin{cases} 0.5 & \text{if } \varphi_c \neq \neg(G_1 \wedge \dots \wedge G_n) \\ 0 & \text{otherwise} \end{cases}$$

**the shorter,
the better**

Evaluation

RQ1 How effective and efficient is our approach to identify boundary conditions in requirement specifications?

RQ2 Is our approach able to identify boundary conditions that cannot be derived by related techniques?








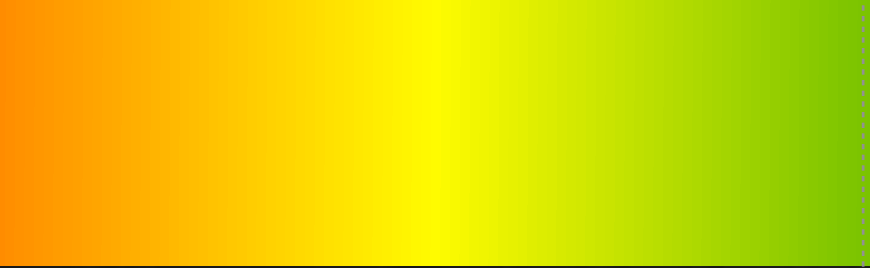

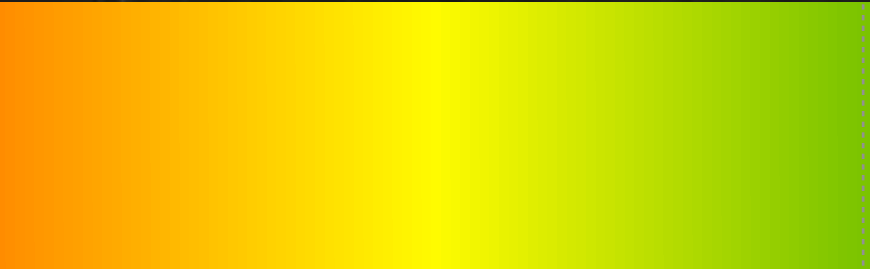

<http://dc.exa.unrc.edu.ar/staff/rdegiovanni/ase2018.html>

The tool: **JGAP**, **LTL2Buchi**, and **aalta** LTL solver.

Evaluation

Case Study	Pattern-based	Tableaux-based	Genetic Algorithm
Achieve-Avoid	1	4	21
Retraction 1	1	1	27
Retraction 2	1	1	22
RailRoadCrossingSystem	-	1	16
MinePump	-	2	18
ATM	-	4	10
Elevator	-	1	7
TCP protocol	-	2	8
Telephone	-	1	24
London Ambulance System	-	1	84
Simple Arbiter	-	TO	15
Prioritized Protocol	-	TO	13
Round Robin Arbiter	-	TO	37
Load Balancer	-	TO	3
Lift Controller	-	TO	3
AMBA	-	TO	2

Summary

	Pattern-based	Tableaux-based	Genetic Algorithm
scalability			
readability			
applicability			
completeness			

Applicability and Usability

Control Synthesis Problem



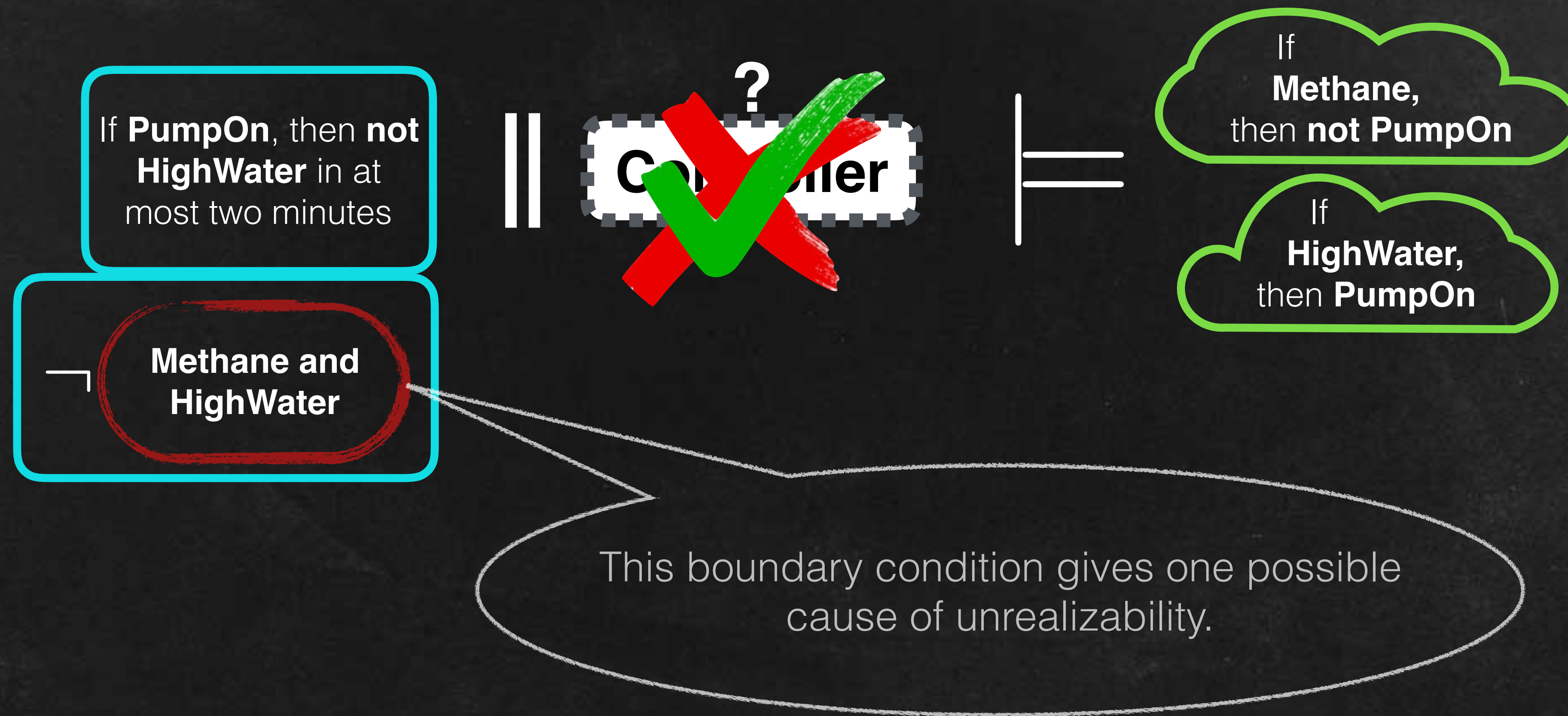
✓ realizable

✗ unrealizable

Can boundary conditions
explain why the specifications are
unrealizable?

Applicability and Usability

Mine Pump Controller



Remarks

- Novel application of genetic algorithms in the context of software engineering
- More general and scalable automated technique for boundary condition computation
- Enables the application of boundary conditions for requirements engineering problems with increased demands of scalability

Questions?

Thanks

Evaluation (with spec sizes)

Case Study	Spec size	Pattern-based	Tableaux-based	Genetic Algorithm
Achieve-Avoid	3	1	4	21
Retraction 1	2	1	1	27
Retraction 2	2	1	1	22
RailRoadCrossingSystem	4	-	1	16
MinePump	3	-	2	18
ATM	3	-	4	10
Elevator	2	-	1	7
TCP protocol	2	-	2	8
Telephone	5	-	1	24
London Ambulance System	5	-	1	84
Simple Arbiter	7	-	TO	15
Prioritized Protocol	7	-	TO	13
Round Robin Arbiter	9	-	TO	37
Load Balancer	11	-	TO	3
Lift Controller	21	-	TO	3
AMBA	27	-	TO	2

Time Comparison (sec.)

Case Study	Pattern-based	Tableaux-based	Genetic Algorithm
Achieve-Avoid	0	2	5
Retraction 1	0	0	17
Retraction 2	0	0	16
RailRoadCrossingSystem	-	1	17
MinePump	-	9	7
ATM	-	10	7
Elevator	-	0	0
TCP protocol	-	1	10
Telephone	-	5	53
London Ambulance System	-	5	8491
Simple Arbiter	-	TO	406
Prioritized Protocol	-	TO	8770
Round Robin Arbiter	-	TO	152
Load Balancer	-	TO	6578
Lift Controller	-	TO	2853
AMBA	-	TO	7541

Genetic Algorithm Configuration

Case Study	Spec size	Pop size	Chrom size	Generations
Achieve-Avoid	3	100	20	50
Retraction 1	2	100	20	50
Retraction 2	2	100	20	50
RailRoadCrossingSyste	4	100	20	50
MinePump	3	100	20	50
ATM	3	100	20	50
Elevator	2	100	20	50
TCP protocol	2	100	20	50
Telephone	5	500	50	50
London Ambulance	5	200	50	50
Simple Arbiter	7	100	50	50
Prioritized Protocol	7	100	50	50
Round Robin Arbiter	9	100	20	50
Load Balancer	11	200	50	50
Lift Controller	21	100	50	50
AMBA	27	100	50	50